

When does a slime mould compute?

Memorial University
24 July 2014

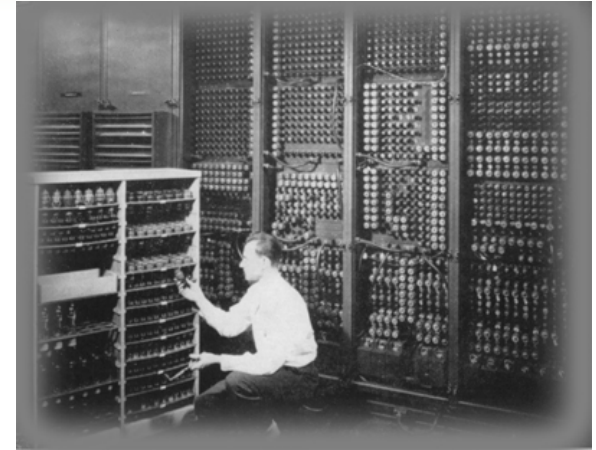
Susan Stepney

Non-Standard Computation Research Group

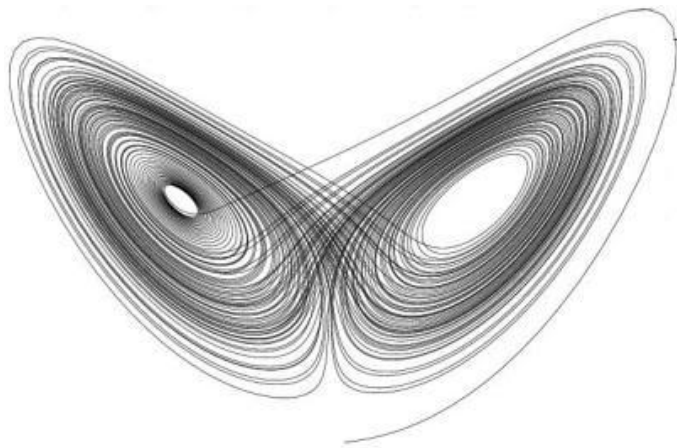
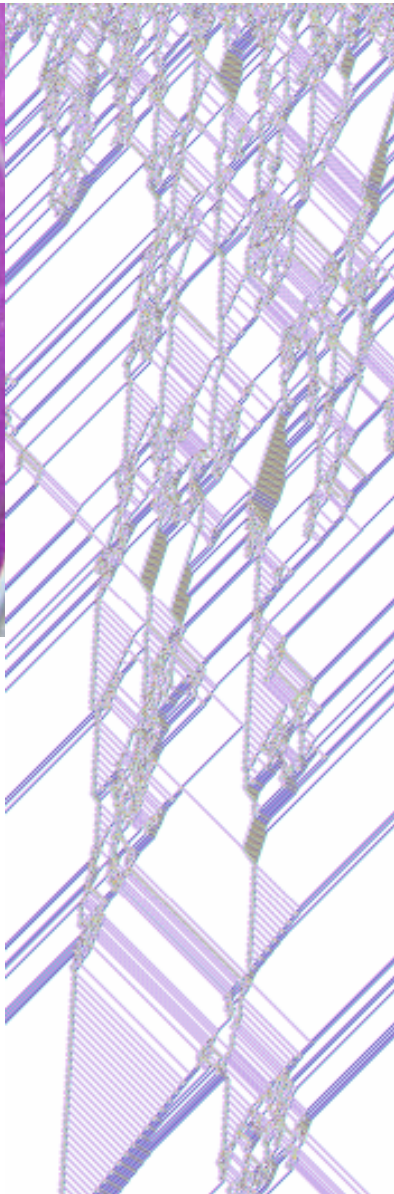
Department of Computer Science



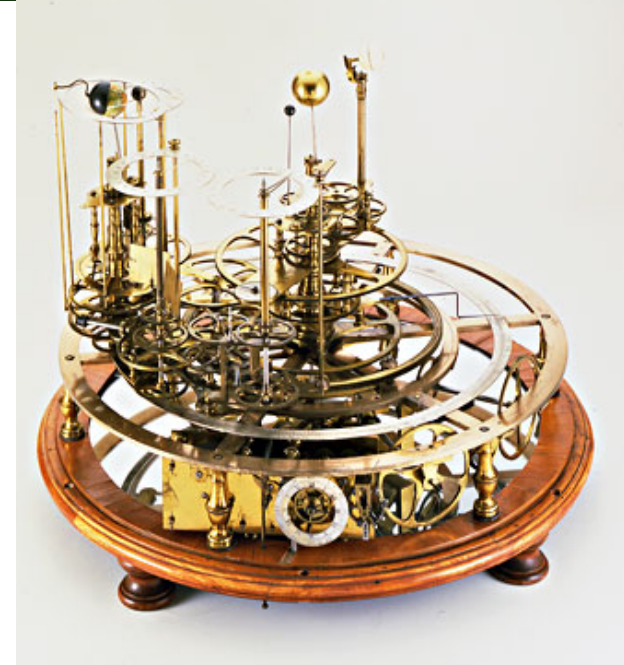
if I say “computer”, you probably think ...



but do you think ...?



you shouldn't be surprised



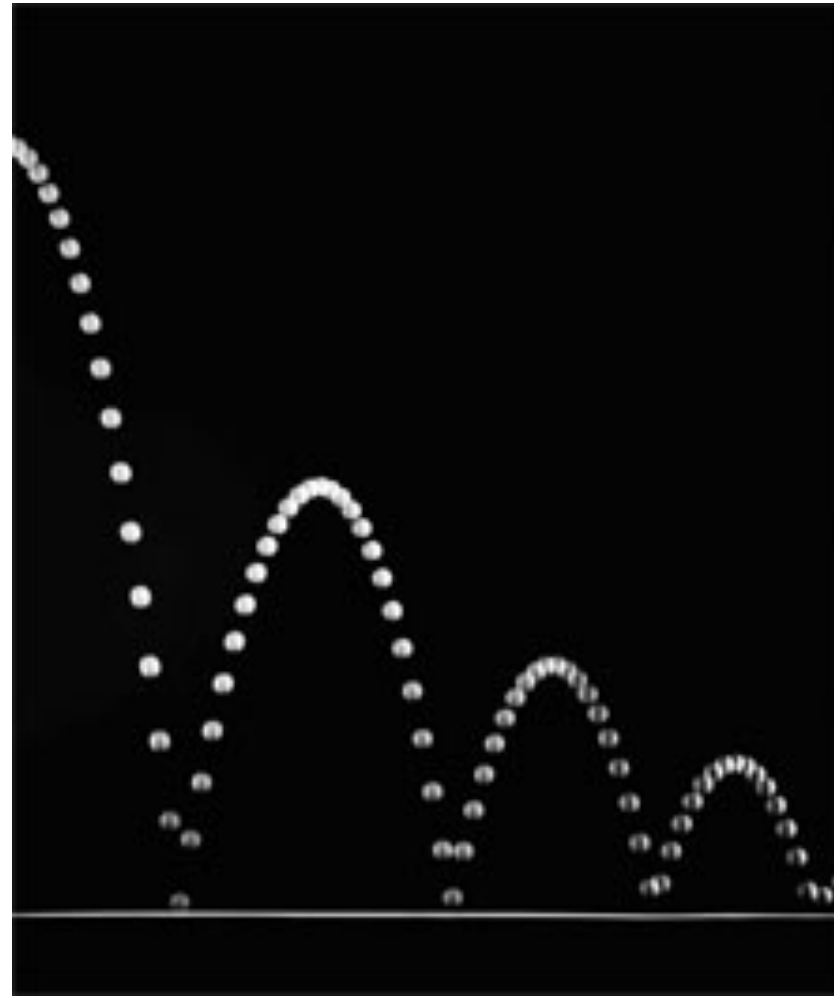
Antikythera, c.80CE

but what about these?



what is a computer?

- What does it mean to say that some physical system is “running” a computation?
 - as opposed to just “doing its thing”
- When does a physical system “compute”?



three steps to computers

first we need to answer:

- what is **science**?
 - how we represent a physical system as an “abstract model”

this will let us answer:

- what is **engineering**?
 - how we instantiate an abstract model as a physical system

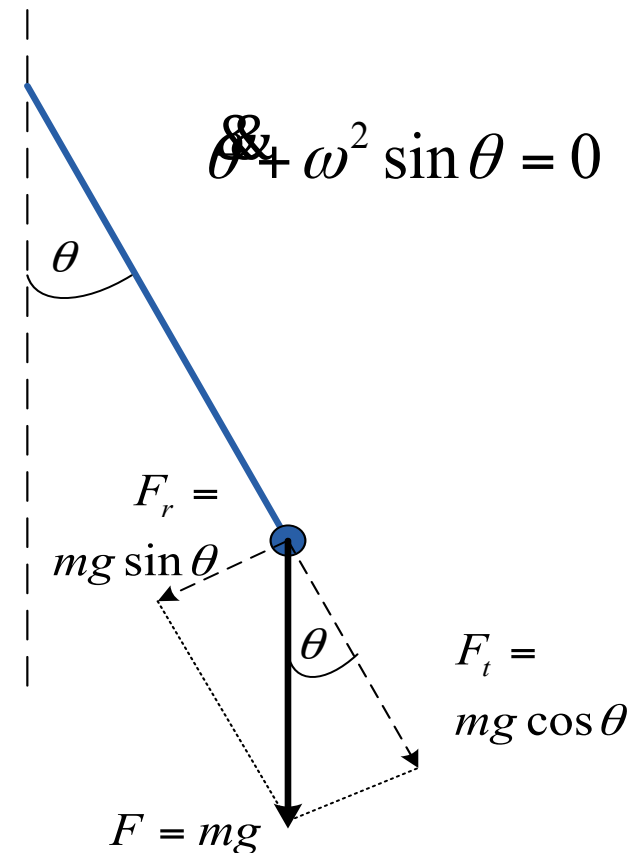
and then this will let us answer:

- what is **computing**?
 - how we instantiate a computational model in a physical system

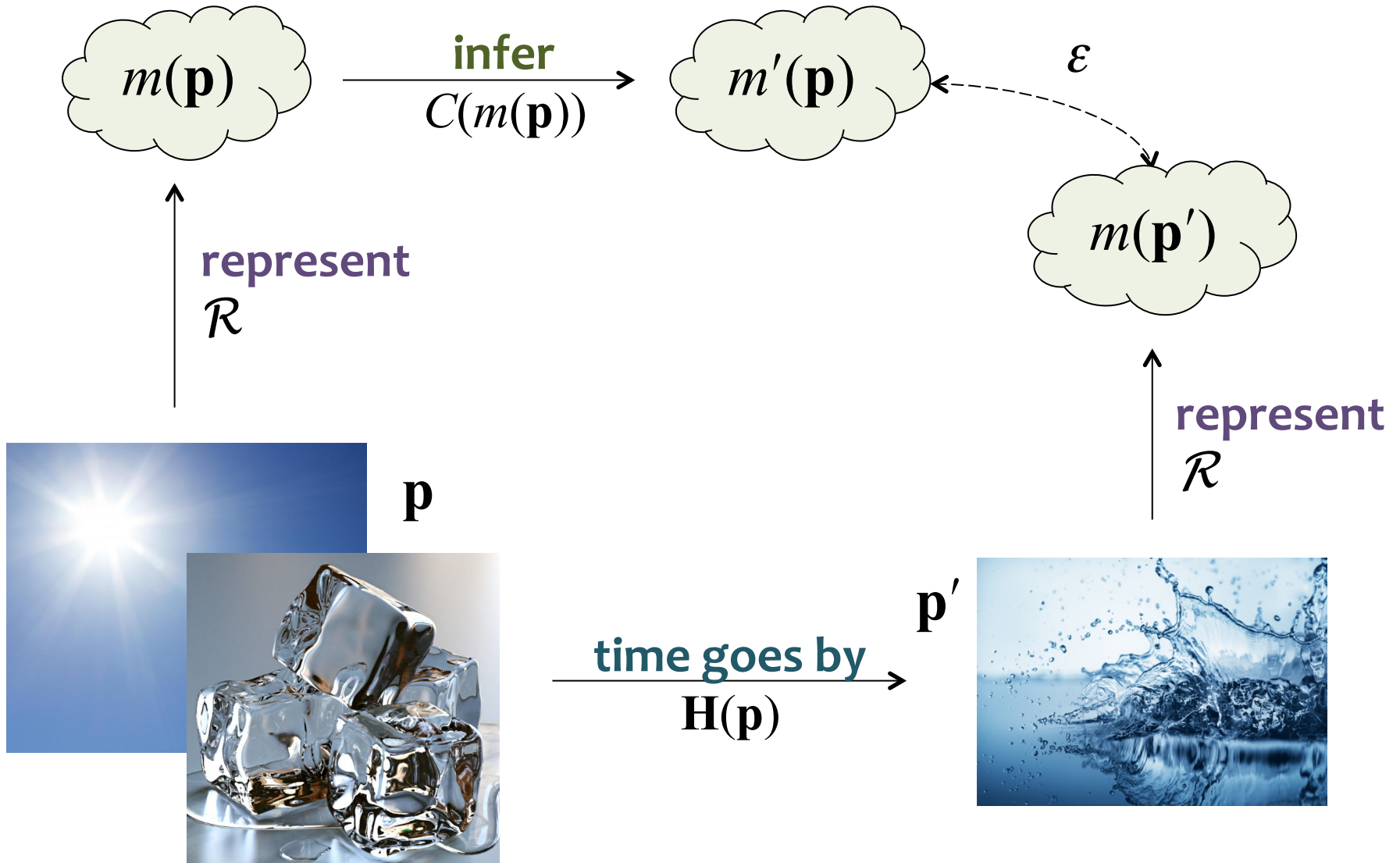


science

physical system $\xrightarrow{\text{represented}}$ abstract model
 $\xleftarrow{\text{instantiated}}$

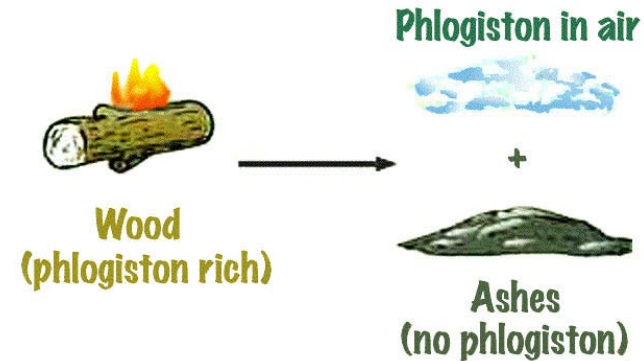


science

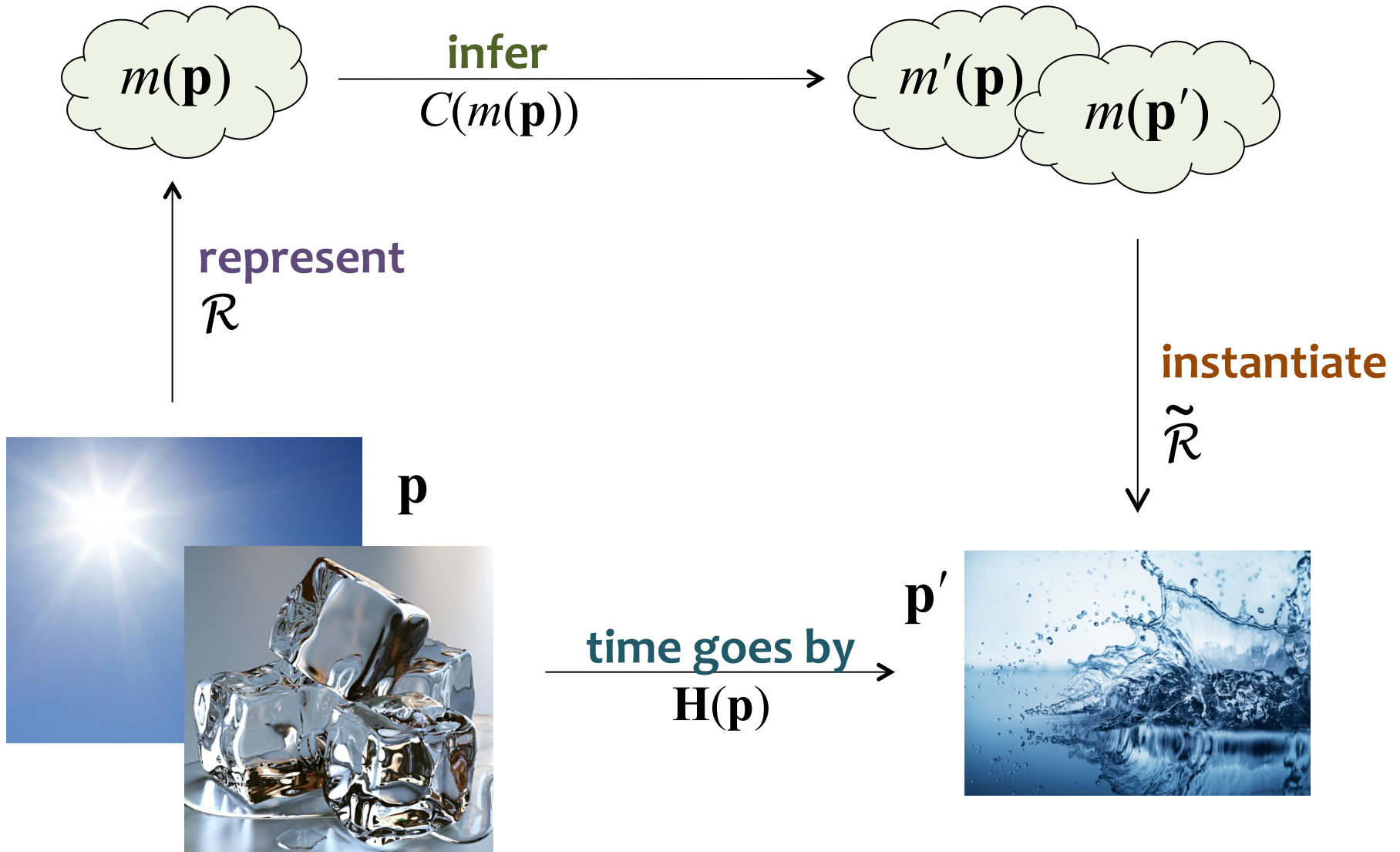


science

- a “good” theory makes ϵ “small enough”
 - among other things...
- if ϵ is too large, change the theory
 - reality trumps theory
- a theory is a *model* of reality
 - models are *always* approximations
 - approximations *break down* outside the model’s valid domain
- a good theory allows prediction *without* needing a “reality check” every time a prediction is made
 - within the domain where the approximations hold

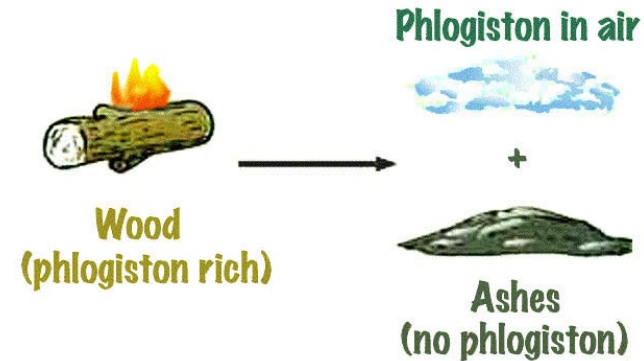


good scientific theory : prediction

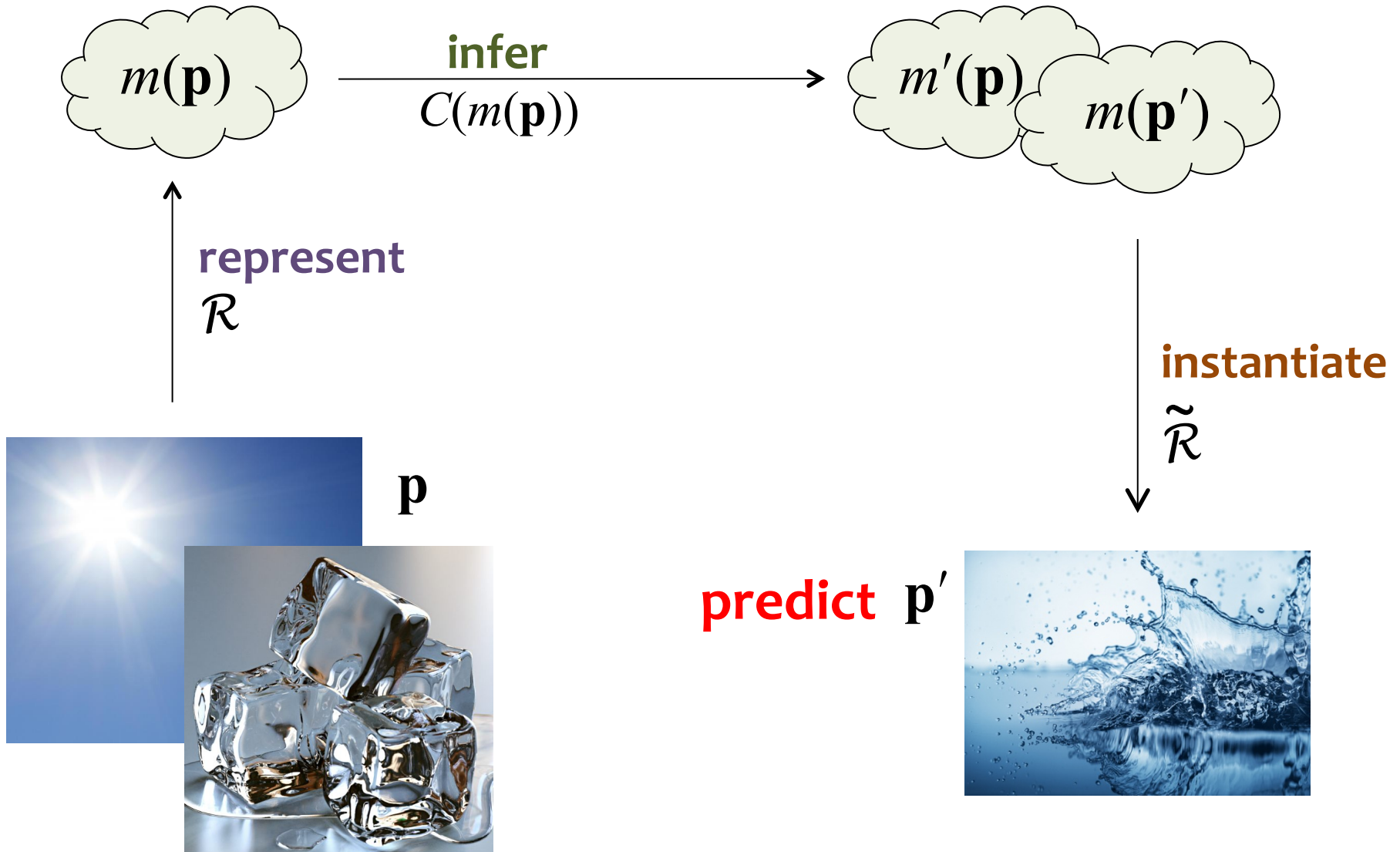


science

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good scientific theory : prediction

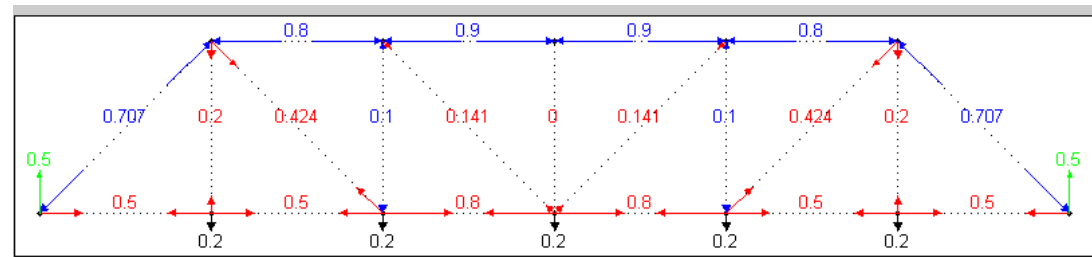


definition

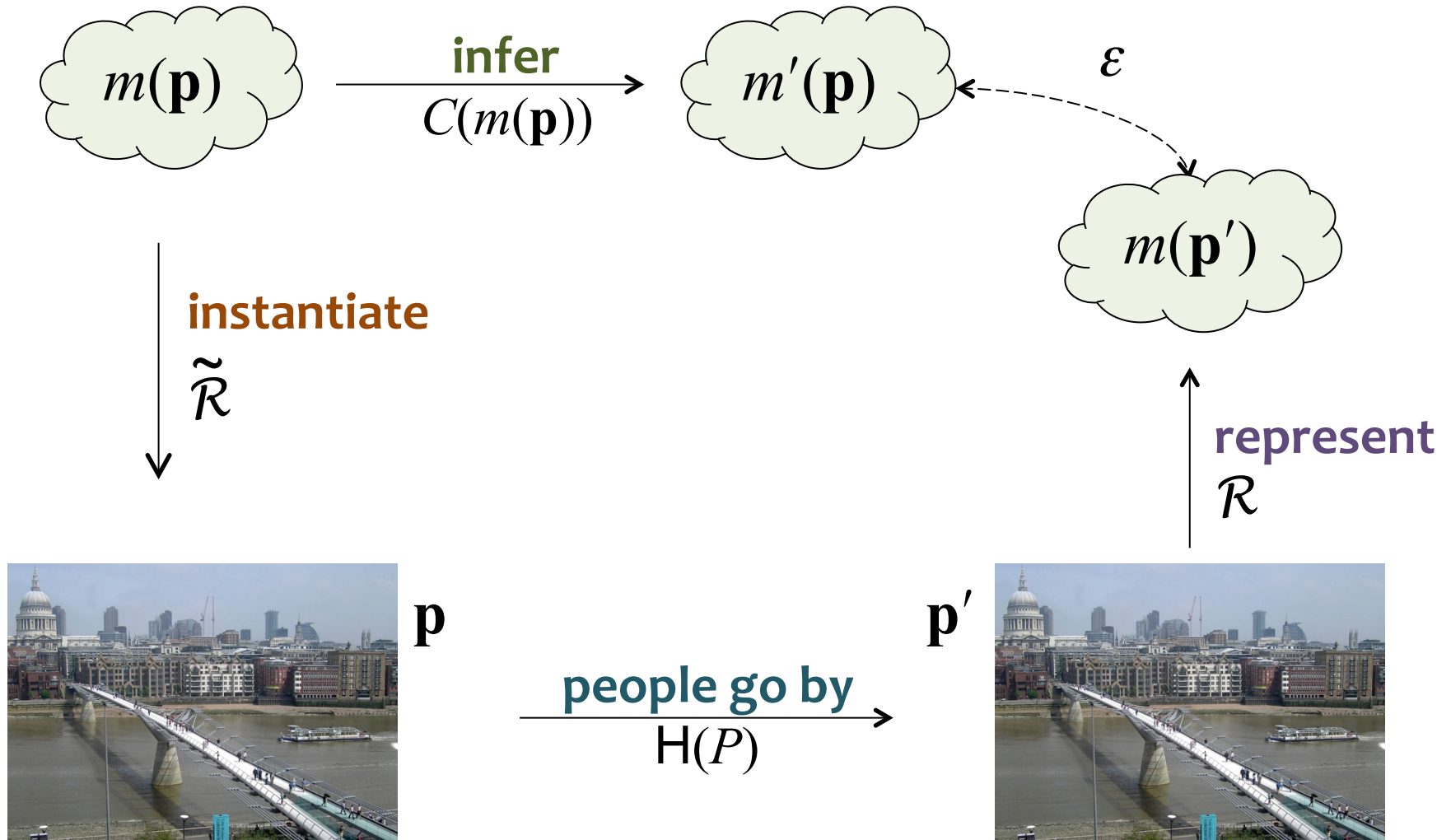
here, prediction is :

using an **abstract** dynamics **C**
of a ***well-characterised*** physical system
to ***infer*** its **physical** dynamics **H**
(subject to a representation \mathcal{R})

technology / engineering

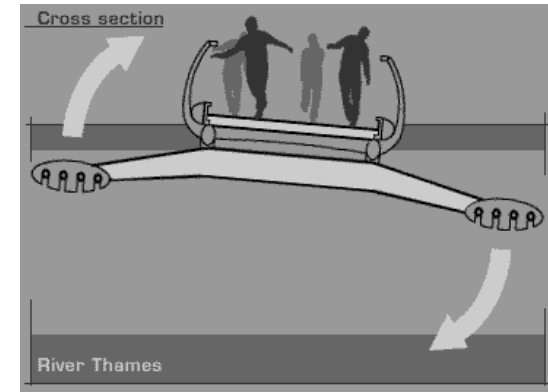


technology / engineering

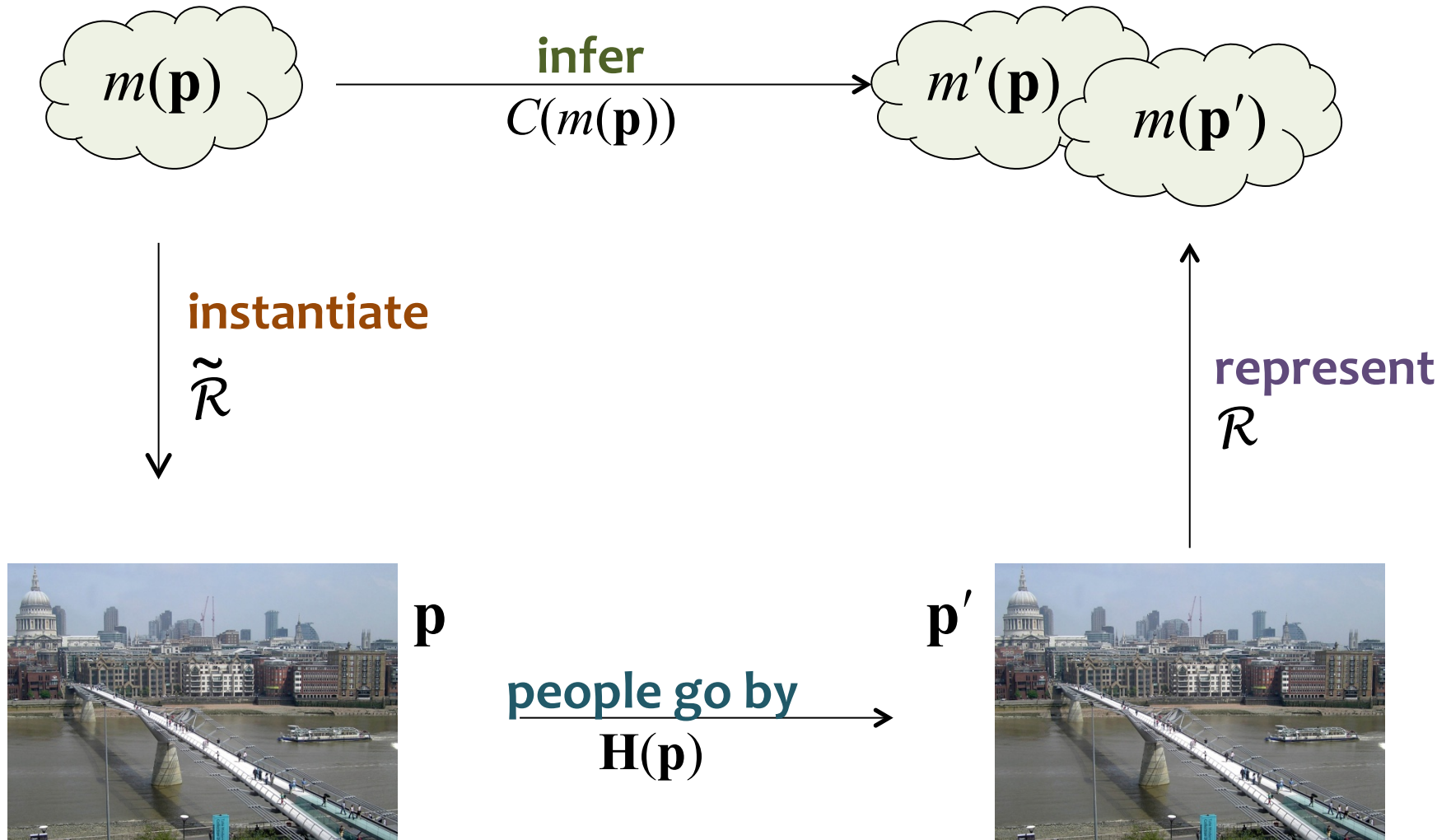


engineering

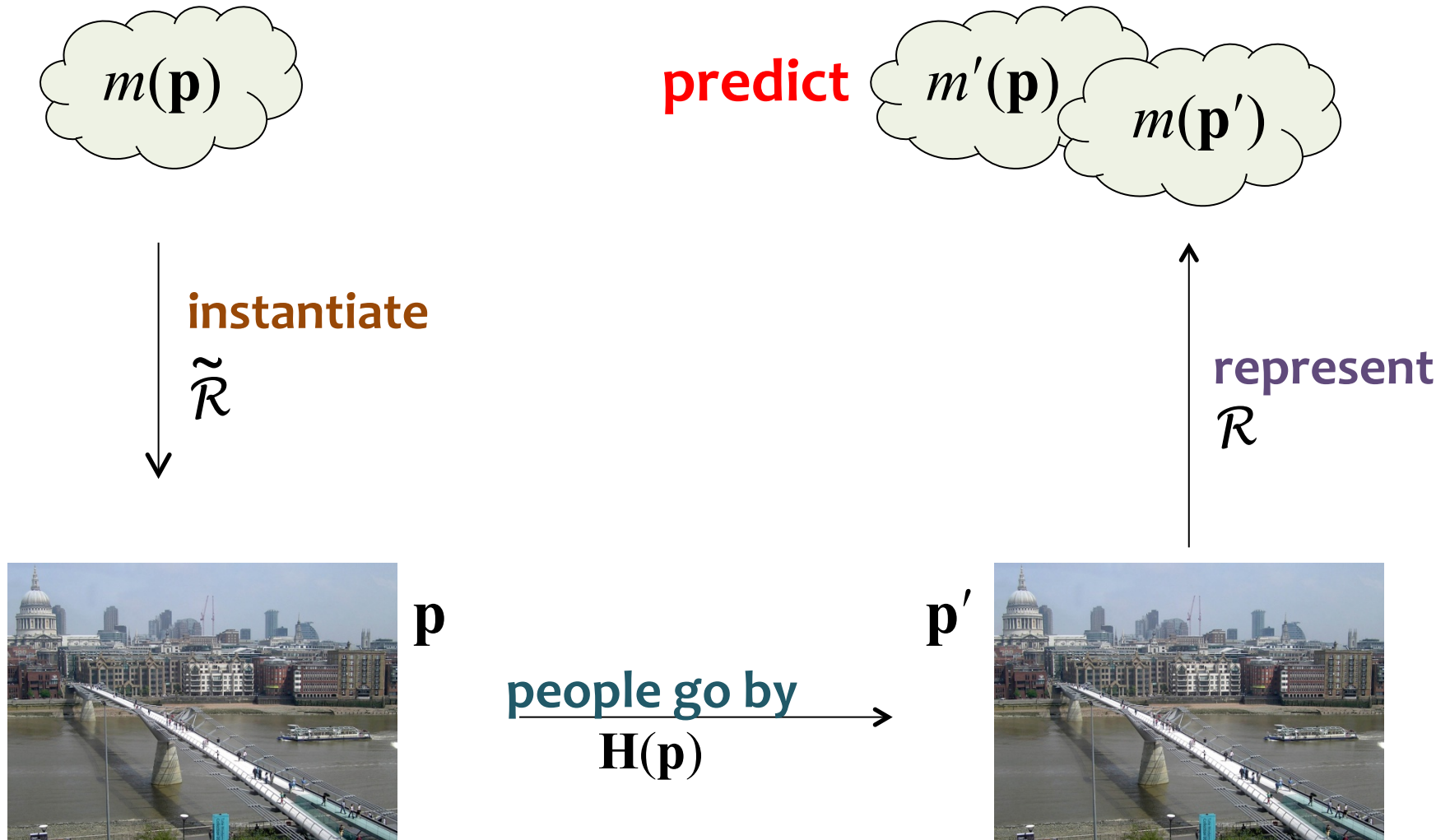
- a “good” instantiation \mathbf{p} makes ε “small enough”
 - among other things...
- if ε is too large, \mathbf{p} needs to be changed
 - the (desired) model trumps reality
- a theory is a *model* of reality
 - models are *always* approximations
 - approximations *break down* outside the model’s valid domain
- a good instantiation allows use *without* needing a “theory check” every time the system is used
 - within the domain where the approximations hold



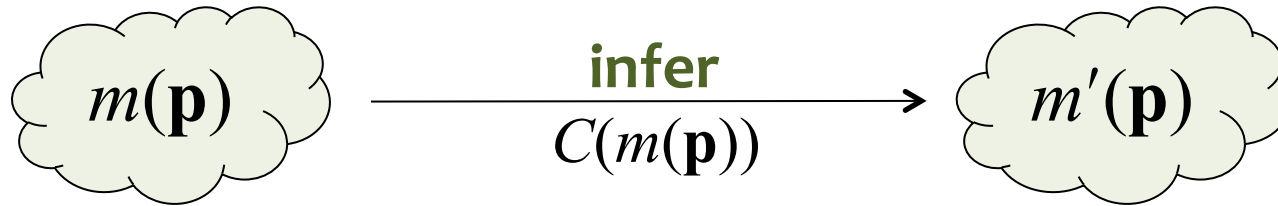
well-engineered technology



well-engineered technology



inferring



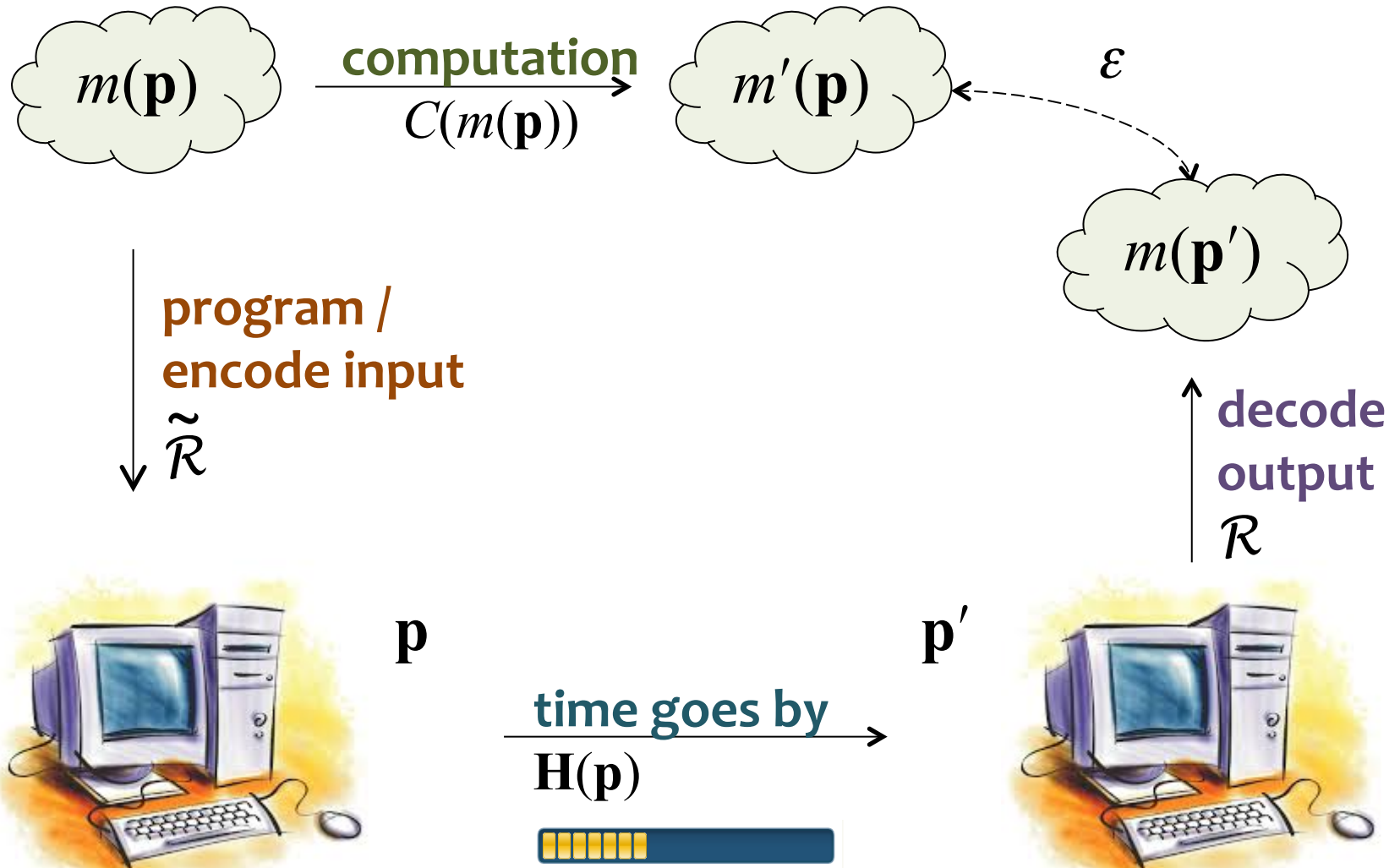
- who, or what, is doing the inferring?
 - a long complicated calculation
 - done with pen and paper ...
 - ... or with a computer!

computing



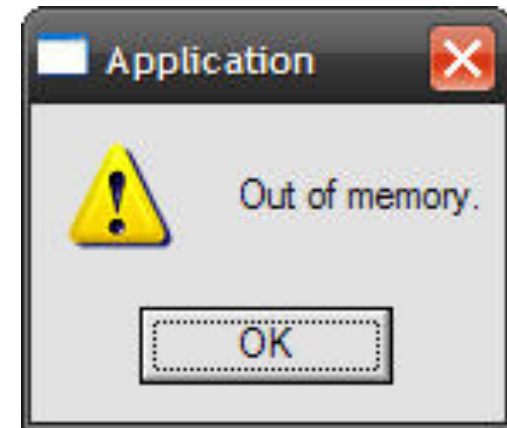
```
1: for  $t = 0$  to  $p - 1$  {each automaton state in the ash} do
2:   while there is a live cell  $c$  that has not been processed yet do
3:     if  $c \notin \Omega$  { $c$  is not contained in any oscillator} then
4:       {create a new oscillator  $O$  containing just  $c$  and its current state}
5:        $O := \{(c, \sigma_t(c))\}; \Omega := \Omega \cup \{O\}$ 
6:     else { $c$  is in an oscillator, with  $(c, S) \in O$ }
7:        $S := S + \sigma_t(c)$  {update  $c$ 's state list with  $c$ 's current state}
8:     end if
9:   for each  $n \in N(c)$  {each of cell  $c$ 's neighbourhood cells} do
10:    if  $\sigma_t(n) = \blacksquare$  { $n$  is alive}
11:    or  $\sigma_t(n) = \square$  and  $|N_{\blacksquare}(n)| \geq 3$ 
12:      { $n$  is dead and has three or more live cells in its neighbourhood} then
13:         $O := O \cup \{(n, \sigma_0(n), \dots, \sigma_t(n))\}$  {add  $n$  to  $O$ }
14:      end if
15:    end for
16:  if any of the cells  $n$  added to  $O$  are already a member of another oscillator  $R$  then
17:     $O := O \cup R; \Omega := \Omega - \{R\}$  {combine  $O$  and  $R$ }
18:  end if
19: continue recursively processing all neighbourhood cells  $n$  added to  $O$ 
20: end while
21: end for
```


computing

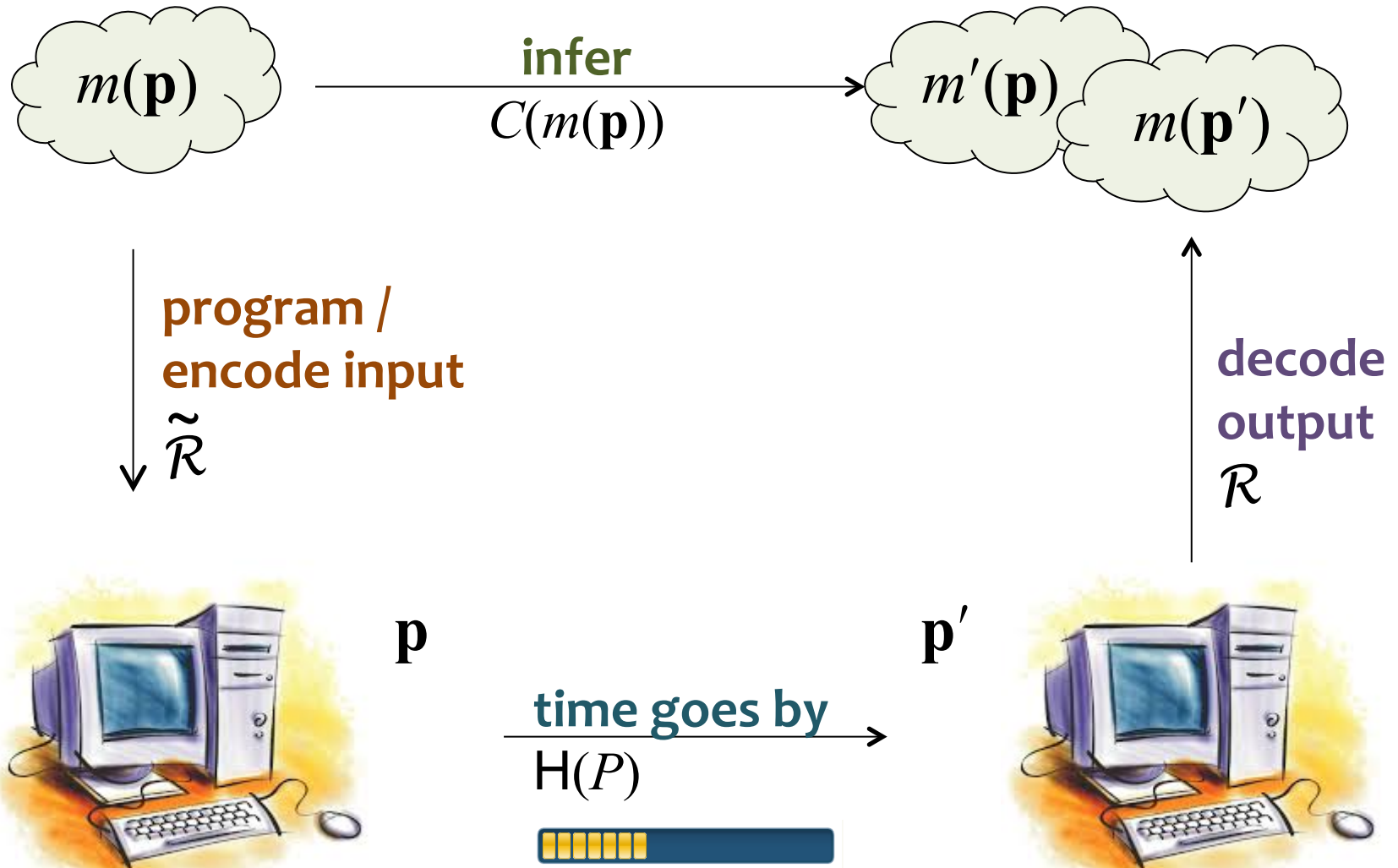


computing

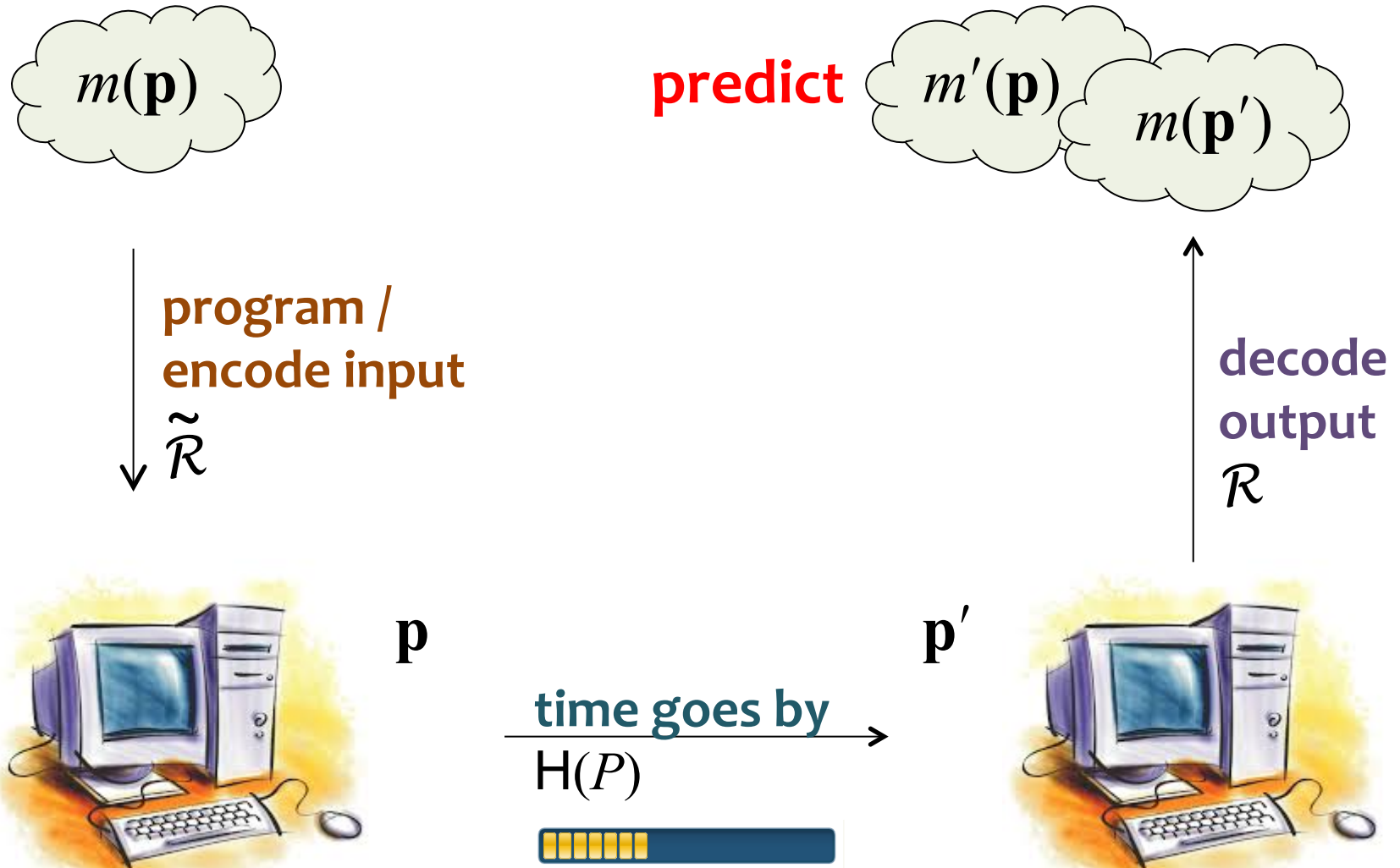
- a “good” instantiated computer p makes ε “small enough”
 - among other things...
- if ε is too large, p needs to be changed
 - the (desired) computation trumps reality
- a theory is a *model* of reality
 - models are *always* approximations
 - approximations *break down* outside the model’s valid domain
- a well-instantiated computer allows use *without* needing a “computation check” every time the system is used
 - within the domain where the approximations hold



well-engineered computer / program



well-engineered computer / program



definition

computation is :

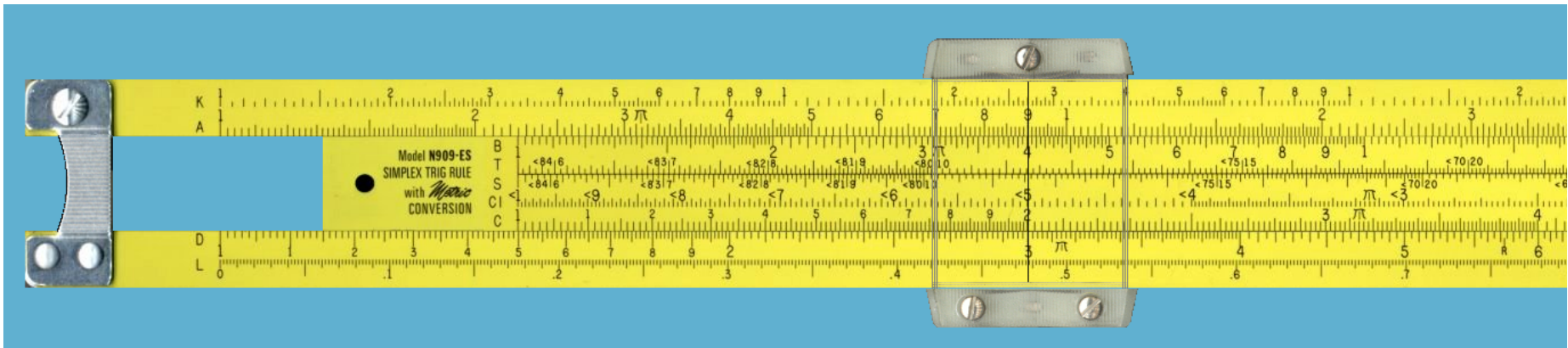
using the **physical** dynamics **H**
of a *well-engineered* physical system
to *predict* an **abstract** dynamics **C**
(subject to an encoding \mathcal{R})

unconventional computing

- there is nothing in the definition about the nature of the physical system
 - beyond being “well-engineered”
- it doesn't have to be silicon
- it doesn't have to be a conventional computer
- we can use this definition to understand how unconventional physical systems compute

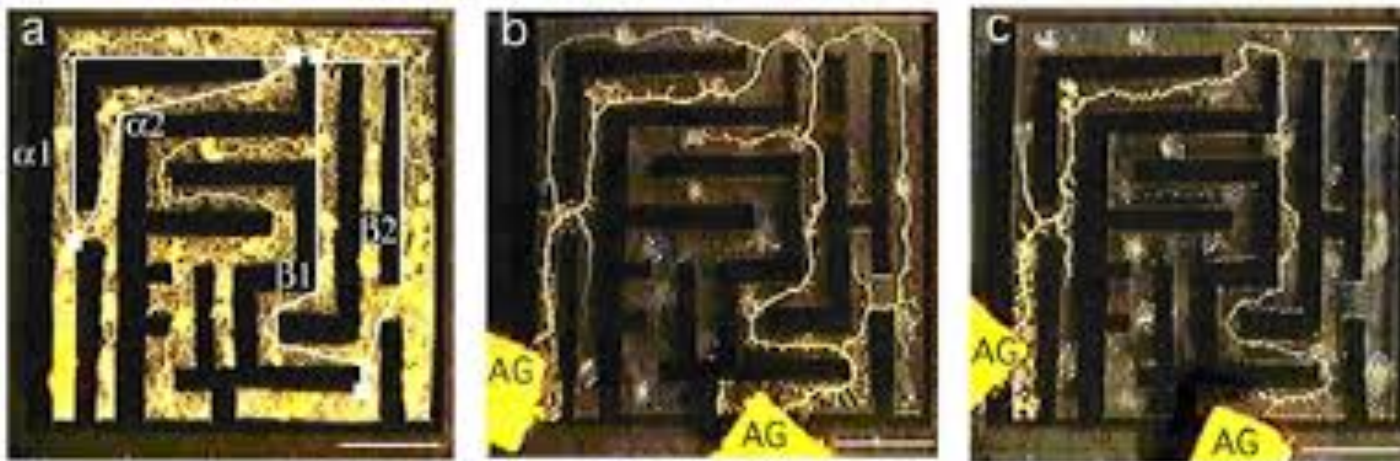
example : wooden sticks

- *calculation* : multiplication
- *theory* : how lengths of bits of wood combine
 - they add together linearly
- *instantiation* : abstract numbers instantiated as physical lengths
- *real world* : join lengths together
- *output* : read off the total length
 - logarithmic scale : so *multiplies* the values



example : slime mould

- *calculation* : solving a maze
- *theory* : how slime moulds behave in presence of food
 - they minimise distances
- *instantiation* : chopped up slime mould covers maze
 - food sources at entrance and exit
- *real world* : slime mould contracts, joining the sources
- *output* : read off path taken by slime mould



requirements for
a physical computer

(i) a well-characterised substrate



$$\frac{\partial U_i^0(\vec{x}, t)}{\partial t} - \frac{1}{Re} \sum_{j=1}^3 \frac{\partial^2 U_i^0(\vec{x}, t)}{\partial x_j \partial x_j} + \frac{\partial U_4^0(\vec{x}, t)}{\partial x_i} + \sum_{j=1}^3 \left(U_j^0(\vec{x}, t) \frac{\partial U_i^0(\vec{x}, t)}{\partial x_j} + \sum_{j_1=1}^4 \int_{\vec{x}_1, t_1} U_{jj_1}^1(\vec{x}, t; \vec{x}_1, t_1) \frac{\partial U_{ij_1}^1(\vec{x}, t; \vec{x}_1, t_1)}{\partial x_j} d\vec{x}_1 dt_1 \right) = 0$$

$$\sum_{j=1}^3 \frac{\partial^2 U_4^0(\vec{x}, t)}{\partial x_j \partial x_j} + \sum_{i,j=1}^3 \left(\frac{\partial U_i^0(\vec{x}, t)}{\partial x_j} \frac{\partial U_j^0(\vec{x}, t)}{\partial x_i} + \sum_{j_1=1}^4 \int_{\vec{x}_1, t_1} \frac{\partial U_{ij_1}^1(\vec{x}, t; \vec{x}_1, t_1)}{\partial x_j} \frac{\partial U_{jj_1}^1(\vec{x}, t; \vec{x}_1, t_1)}{\partial x_i} d\vec{x}_1 dt_1 \right) = 0$$

$$\frac{\partial U_{ij_1}^1(\vec{x}, t; \vec{x}_1, t_1)}{\partial t} - \frac{1}{Re} \sum_{j=1}^3 \frac{\partial^2 U_{ij_1}^1(\vec{x}, t; \vec{x}_1, t_1)}{\partial x_j \partial x_j} + \frac{\partial U_{ij_1}^1(\vec{x}, t; \vec{x}_1, t_1)}{\partial x_i} + \sum_{j=1}^3 \left(U_j^0(\vec{x}, t) \frac{\partial U_{ij_1}^1(\vec{x}, t; \vec{x}_1, t_1)}{\partial x_j} + U_{jj_1}^1(\vec{x}, t; \vec{x}_1, t_1) \frac{\partial U_i^0(\vec{x}, t)}{\partial x_j} \right) = 0$$

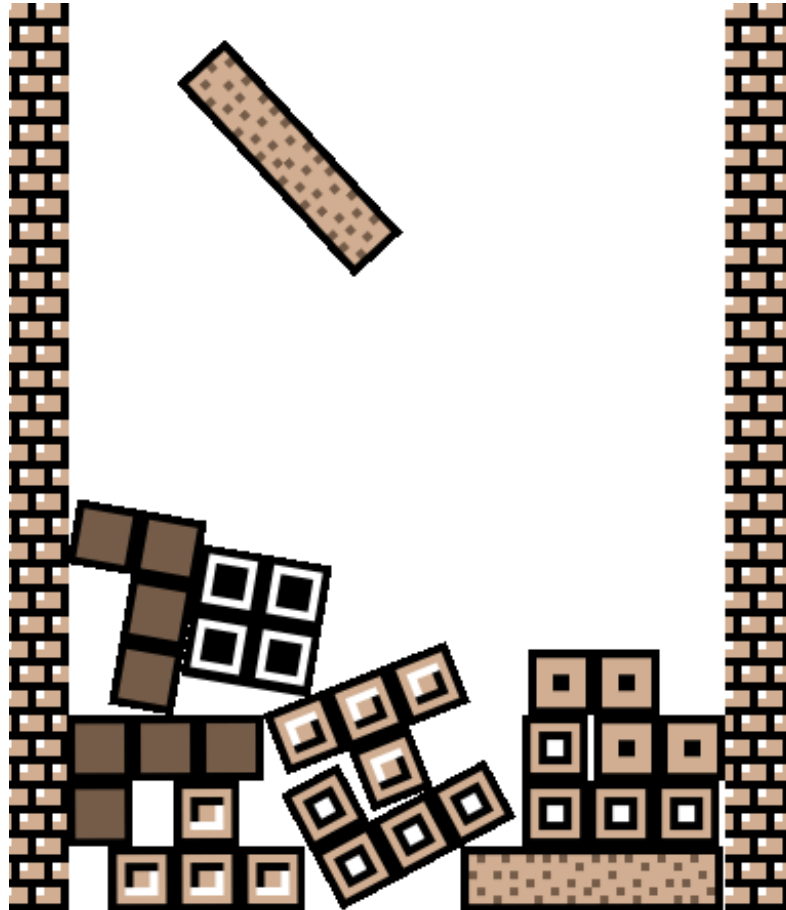
$$\sum_{j=1}^3 \frac{\partial^2 U_{4j_1}^1(\vec{x}, t; \vec{x}_1, t_1)}{\partial x_j \partial x_j} + \sum_{i,j=1}^3 2 \frac{\partial U_i^0(\vec{x}, t)}{\partial x_j} \frac{\partial U_{jj_1}^1(\vec{x}, t; \vec{x}_1, t_1)}{\partial x_i} = 0$$

- including domain of applicability
 - eg, “shortest path” is a rough approximation, for small systems

substrate theories

- well-developed
 - solid state transistors
 - classical mechanics, quantum mechanics
 - reaction-diffusion chemistry
- phenomenological
 - biology
 - ♦ extrapolation and scaling issues
- naïve
 - approximate
 - ♦ shortest path
 - counterfactual
 - ♦ unbounded speeds
 - ♦ non-atomic

(ii) a well-engineered instantiation



engineering issues

- theory composition
 - multiple components
 - multiple kinds of components
 - interconnections
 - control
 - *programming*
- scaling
 - interpolation
 - extrapolation
 - ◆ model breaks down

(iii) a pre-defined encoding/decoding



Nope



(iv) and a natural fit to the problem



not here ...



HM Revenue
& Customs

Tax Return 2010

Tax year 6 April 2009 to 5 April 2010

natural fit

- the fit between the desired abstract dynamics and the possible physical dynamics
- small “semantic gap”
 - actually pretty poor for conventional computers!
 - “torturing” silicon to implement boolean logic
- smaller gap with other substrates, other computational models?
 - analogue computers
 - other unconventional approaches

acknowledgments

Clare Horsman, Susan Stepney, Rob C. Wagner, Viv Kendon.

When does a physical system compute?

Proceedings of the Royal Society A, **470**(2169):20140182

doi: 10.1098/rspa.2014.0182